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It has a least el't 0
which is followed by next least el't 1
" " " " " " " " " " " "

$\mu: \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet$
 $\quad \quad \quad \circ^{\mu} \quad 1^{\mu} \quad 1^{\mu} + 1^{\mu}$

What claim says ω there is $M \models Th(N)$ with an "infinite element" (!)

μ

Once there is one infinite element,
there are many.

$$M = \underbrace{\quad \quad \quad}_{\mathbb{N}} \quad \quad \quad q-1 \quad q \quad q+1$$

So: even if we take our "axioms" to be all first-order statements which are true of $(\mathbb{N}, +, \cdot, <, 0, 1)$... these axioms not enough to characterize $\mathbb{N}$!

(SP)

PF of claim: - Let $L^* = L \cup \{c\}$ where

c is a new constant symbol

- Let φ_n be $\underbrace{1+1+\dots+1}_{n\text{-times}} < c$

- Let $T' = Th(N) \cup \{\varphi_1, \varphi_2, \dots\}$

- If $\Delta \subseteq T' \cup \Gamma_{\text{th}(N)}$, then

$\Delta \subseteq Th(N) \cup \{\varphi_1, \dots, \varphi_N\}$ for some large enough N .

- Let $\mathcal{M}$ be L^* -expansion of N where we let $c^{\mathcal{M}} = N+1$

- Then $\mathcal{M} \models \Delta$

- Thus $T' \cup \Gamma_{\text{th}(N)}$ is sat.

$\Rightarrow$ there is $\mathcal{M} \models T'$ by compactness

- In $\mathcal{M}$, $c^{\mathcal{M}} > n = 1+1+\dots+1$ for all n ✓

Lowenheim-Skolem Theorem

Recall: If $\mathcal{M}, \mathcal{N}$ are L -structs

write $\mathcal{M} \subseteq \mathcal{N}$ ("M a substructure of N")

If $\mathcal{M} \subseteq \mathcal{N}$, and $f^{\mathcal{M}} = f^{\mathcal{N}}|_{\mathcal{M}}$

$R^{\mathcal{M}} = R^{\mathcal{N}}|_{\mathcal{M}}$

$c^{\mathcal{M}} = c^{\mathcal{N}}$ for all $f/R/c$ symbols.

Equivalency: (why?)

$\mathcal{M} \subseteq \mathcal{N}$ if

$\mathcal{M} \models \varphi(\vec{a}) \iff \mathcal{N} \models \varphi(\vec{a})$

for all atomic formulas φ and $\vec{a} \in \mathcal{M}$

- write $M < N$ ("M is elementary substructure of N" "N is elem extension of M")
 $(\models M \models \text{N and } M \models \varphi(\vec{a}) \text{ iff } N \models \varphi(\vec{a})$
 for all formulas $\varphi(\vec{v})$ and $\vec{a} \in M$.

- an injective $j: M \rightarrow N$ is an embedding of M into N if $j[M] \subseteq N$
 - is an elementary embedding if $j[M] < N$, i.e. for all φ and $\vec{a} \in M$
 $M \models \varphi(\vec{a}) \text{ iff } N \models \varphi(j(\vec{a}))$

~~on the other hand~~

- structures strongly resemble their elem substructures - if $M < N$ then e.g.
 every sentence φ true in N iff in M
 - yet we'll prove: given an infinite M , we can find elem extensions of M as big as we want, and elem substructures that are as small as possible.

Def'n Spcs M on $\mathcal{L}$ -structure. let $\mathcal{L}_M$ be lang. obtained by adding to $\mathcal{L}$ a constant symbol m for each el't of M .

(e.g. if $M = (N, +, \cdot, <, 0, 1)$ then $\mathcal{L}_M = \{+, \cdot, <, 0, 1, 2, 3, \dots\}$)

(60)

The atomic diagram of M (write: $\text{Diag}(M)$)
 is the theory $\{\varphi(m_1, \dots, m_n) : \varphi \text{ an atomic or neg'n of atomic formula and } M \models \varphi(m_1, \dots, m_n)\}$

(viewing M as an $\mathcal{L}_M$ -structure here, interpreting every new constant symbol as the el't it represents)

The elementary diagram (write: $\text{Diag}_{el}(M)$)
 is $\{\varphi(\vec{m}) : \varphi(\vec{v}) \text{ an } \mathcal{L}\text{-formula and } M \models \varphi(\vec{m})\}$

Lemma Spc $\mathcal{K}$ is an $\mathcal{L}_M$ -structure

① if $\mathcal{K} \models \text{Diag}(M)$, then (viewing $\mathcal{K}$ as an $\mathcal{L}$ -structure) there is an embedding of M into $\mathcal{K}$

② if $\mathcal{K} \models \text{Diag}_{el}(M)$ then there is an elementary embedding of M into $\mathcal{K}$.

PF: ① Define $j: M \rightarrow \mathcal{K}$ by $j(m) = m^{\mathcal{K}}$
 i.e. we send m to interp (in $\mathcal{K}$) of constant symbol for m .

— j is "clearly" an embedding. Explicitly:

— given distinct $m_1, m_2 \in M$, we have $m_1 \neq m_2 \in \text{Diag}(M)$

— thus $j(m_1) = m_1^{\mathcal{K}} \neq m_2^{\mathcal{K}} = j(m_2)$

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- so j is injective
- given constant symbol $c \in \mathcal{L}$ have $c^{\mathcal{M}} = m$ for some $m \in M$, so $c = m \in \text{Dom}(\text{Diag}(\mathcal{M}))$, so $j(c^{\mathcal{M}}) = j(m) = m^{\mathcal{N}} = c^{\mathcal{N}}$
- given $R \in \mathcal{L}$, if $R^{\mathcal{M}}(m_1, \dots, m_n)$ then $R(m_1, \dots, m_n) \in \text{Diag}(\mathcal{M})$, so $R^{\mathcal{N}}(j(m_1), \dots, j(m_n))$
- given $f \in \mathcal{L}$, if $f^{\mathcal{M}}(m_1, \dots, m_n) = m$ then $f(m_1, \dots, m_n) = m \in \text{Diag}(\mathcal{M})$ so $f^{\mathcal{N}}(j(m_1), \dots, j(m_n)) = j(m)$
- $\Rightarrow j$ an embedding ✓

② observe in this case the map j is elementary. ✓

if $\mathcal{M}$ is an infinite struct. then Lemma + compactness allows us to build elem extensions of $\mathcal{M}$ that are as big as we want.

Theorem (Upward Lowenheim-Skolem)
 Let $\mathcal{M}$ be an infinite $\mathcal{L}$ -structure, and fix a cardinal $\kappa \geq |\mathcal{M}|$. Then there exists an $\mathcal{L}$ -structure $\mathcal{N}$ s.t. $|\mathcal{N}| = \kappa$ and there is an elementary embedding $j: \mathcal{M} \rightarrow \mathcal{N}$.

(62)

PF. - fix $k \geq |M|$

- we assume $|L| \leq |M|$ so we can apply version of compactness we proved, but this is unnecessary if one examines proof
- introduce new constants $c_i : i < k$
- for every (i, j) , add sentence $c_i \neq c_j$ to $\text{Diag}_{e_1}(M)$
- let $T = \text{Diag}_{e_1}(M) \cup \{c_i \neq c_j : i, j < k, (i, j) \neq (j, i)\}$
be resulting theory in lang $L' = L \cup \{c_i : i < k\}$
- if $\Delta \subseteq T$ is finite, then contains only finitely many sentences of the form $c_i \neq c_j$.
- if we interpret symbols c_i appearing in Δ as arbitrary (but distinct) elts of M ,
get $M \models \Delta$. ↑
finite
- so T is fin. sat. hence satisfiable by compactness.
- since $k = \max\{|M|, |L|\}$ there is $M \models T$ of size $\leq k$
- must have $|M| = k$ since the c_i 's are k -many distinct elts of M
- Since $M \models \text{Diag}_{e_1}(M)$ there is
 $\exists : M \rightarrow N$ elementary ✓
- if we identify el w/ $\exists[M]$, then says $\exists$ elem extension of M of every cardinality

(68)

- "Upward" = larger.
- Can we get smaller? Need following:

Prop'n (Tarski-Vaught test) Spcs $M \subseteq N$.
Then $M < N$ iff we have:

(*) For any formula $\varphi(u, \vec{v})$ and $\vec{a} \in M$
if there is $b \in N$ s.t. $N \models \varphi(b, \vec{a})$ then can
find $c \in M$ s.t. $M \models \varphi(c, \vec{a})$

"If there is a witness for φ in N , can find
a witness in M "

But notice: witness ... from pow. of N !

PF ($\Rightarrow$) if $M < N$, then (*) follows immediately
from elementarity

($\Leftarrow$) Assume (*).

Must check ~~for every~~ for every $\varphi(\vec{v})$ and $\vec{a} \in M$:

$$M \models \varphi(\vec{a}) \text{ (iff } N \models \varphi(\vec{a}))$$



by induction on $\varphi(\vec{v})$

- if $\varphi(\vec{v})$ atomic, get this just from $M \subseteq N$
- if true for φ , also for $\neg \varphi$ since
 $M \models \neg \varphi(\vec{a})$ iff $M \not\models \varphi(\vec{a})$ (iff (inductively) $N \not\models \varphi(\vec{a})$ iff $N \models \neg \varphi(\vec{a})$)
- if true for φ, χ then similarly
easily true for $\varphi \wedge \chi$

(64)

- Sps claim true for $\varphi(w, \vec{v})$, consider $\exists w \varphi(w, \vec{v})$
- If $M \models \exists w \varphi(w, \vec{a})$ then $\exists c \in M$ s.t.
 $M \models \varphi(c, \vec{a})$, thus $N \models \varphi(c, \vec{a})$
 by induction
 so $N \models \exists w \varphi(w, \vec{a})$
- and if $N \models \exists w \varphi(w, \vec{a})$ there is $b \in N$ s.t.
 $N \models \varphi(b, \vec{a})$ so by (*) there is
 $c \in M$ s.t. ~~that~~ $N \models \varphi(c, \vec{a})$
 so by induction $M \models \varphi(c, \vec{a})$ and $M \models \exists w \varphi(w, \vec{a})$
 $\Rightarrow$ Claim holds for all $\varphi(\vec{v})$ and $\vec{a} \in M$ by induction

Prop'n suggests idea:

Given N , to find small-as-possible $M \preceq N$

- find set of witnesses M_1 for all existential
 rules $\exists w \varphi(w)$
- then witnesses M_2 for all $\exists w \varphi(w, \vec{a})$
 with $\vec{a} \in M_1$,
- then witnesses $M_3 \dots \exists w \varphi(w, \vec{a})$
 $\dots \vec{a} \in M_2$
 $\dots$ etc.

then take union to get M .

- This works! But instead: we introduce
 functions that will do this work for us
- somewhat analogous to witnesses property

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"bisp" $\rightarrow$ Def'n An $\mathcal{L}$ -theory T has built-in Skolem functions if for every $\mathcal{L}$ -formula $\varphi(w, \vec{v})$ there is a symbol f s.t.

$$T \models \forall \vec{v} (\exists w \varphi(w, \vec{v}) \Rightarrow \varphi(f(\vec{v}), \vec{v}))$$

Lemma Spcs T is an $\mathcal{L}$ -theory. Then there is $\mathcal{L}^* \supseteq \mathcal{L}$ and $T^* \supseteq T$ a $\mathcal{L}^*$ -theory s.t. T^* has bwp and $|\mathcal{L}^*| = \max(|\mathcal{L}|, \aleph_0)$

Moreover, if $M \models T$ there is an $\mathcal{L}^*$ -expansion $M^* \models T^*$.

Pf. - We build a sequence of languages $\mathcal{L} = \mathcal{L}_0 \subseteq \mathcal{L}_1 \subseteq \mathcal{L}_2 \subseteq \dots$

and theories $T = T_0 \subseteq T_1 \subseteq T_2 \subseteq \dots$
 s.t. T_i is $\mathcal{L}_i$ -theory

- Given $\mathcal{L}_i$, for every $\mathcal{L}_i$ -formula $\varphi(v, \vec{v}, w)$ we introduce an n -ary f'n symbol f_φ
- let $\mathcal{L}_{i+1} = \mathcal{L}_i \cup \{f_\varphi : \varphi(\vec{v}, w) \text{ is } \mathcal{L}_i\text{-formula}\}$

reversing
 $\vec{v}$ and w
 here for
 the reason
 xx
)

- for each $\varphi(\vec{v}, w)$ on $\mathcal{L}_i$ -formula, let Φ_φ be the sentence

$$\forall \vec{v} (\exists w \varphi(\vec{v}, w) \Rightarrow \varphi(\vec{v}, f_\varphi(\vec{v})))$$

- let $T_{i+1} = T_i \cup \{\Phi_\varphi\}$
- let $\mathcal{L}^* = \bigcup_i \mathcal{L}_i$, $T^* = \bigcup_i T_i$