

Due: Monday November 1, noon.

1. Let  $G$  be a bipartite graph in which every vertex has degree at most  $k$ . Show that  $G$  is a subgraph of a bipartite graph in which every vertex has degree  $k$ . Deduce that  $G$  can be partitioned into  $k$  matchings.
2. (i) Say that vertices  $u, v$  in  $G$  are  $k$ -linked if there are  $k$  mutually edge-disjoint paths from  $u$  to  $v$ . Suppose  $u, v, w$  are distinct,  $u, v$  are  $k$ -linked and  $v, w$  are  $k$ -linked. Does it follow that  $u, w$  are  $k$ -linked?  
(ii) Say subsets  $X, Y \subset V(G)$  are  $k$ -joined if  $|X| = |Y| = k$  and there are  $k$  mutually vertex-disjoint paths from  $X$  to  $Y$ . Suppose  $X, Y, Z \subset V(G)$ ,  $X, Y$  are  $k$ -joined and  $Y, Z$  are  $k$ -joined. Does it follow that  $X, Z$  are  $k$ -joined?
3. Let  $G$  be a transportation network with source  $s$  and sink  $t$ . Let  $f$  be a flow with value  $v$  and suppose there is some flow with value larger than  $v$ . Must there be a flow  $f'$  with value larger than  $v$  so that  $f'(e) \geq f(e)$  for every edge  $e$ ?
4. Let  $T$  be a tree. Show that  $T$  has a perfect matching<sup>1</sup> if and only if for every vertex  $v$  exactly one component of  $T \setminus v$  has an odd number of vertices.
5. Let  $G$  be a bipartite graph with parts  $X$  and  $Y$ . Suppose that for each  $v \in V(G)$  there is a total order  $<_v$  on its neighbourhood  $N(v)$ . We say that a matching  $M$  in  $G$  is stable if whenever we have  $x \in X$ ,  $y \in Y$  with  $xy \notin M$  then either  $xy' \in M$  for some  $y' >_x y$  or  $x'y \in M$  for some  $x' >_y x$ . Show that there is a stable matching.

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<sup>1</sup>i.e. a matching covering all vertices